

Name: _____

Problem 1 (18%). Assemble the following machine instructions into **binary**, use spaces to separate fields, and **registers** in their symbolic form (\$ra NOT \$31). Assume absolute jump addresses.

Field 1	Fields 2 and etc	instruction
100011	00011 11111 00000000000000001100	lw \$ra, 12(\$v1)
001000	10010 11001 0000000011111010	addi \$t9, \$s2, 0xFA
000000	00000 11101 00100 11100 000000	sll \$a0,\$sp,0x1C
000000	01011 00000 00101 00000 100000	move \$a1,\$t3
000011	0000000000000000000000000010100	jal \$20
000000	1001001000010000000000000010000	add \$t0,\$s2,\$t0

Problem 2 (7%).

Assume each part is **independent**. Assume absolute jump & branch addresses (no pc relative). Fill in only registers that changed!

What is the value of the register or memory contents **after** the execution of the instruction.

Assume pc = 1000; \$t1=12; \$t2=3; \$ra=50; memory[12]=0x17761453;

instruction	pc	\$ra	\$t1	memory[12]
lw \$ra, 12(\$0)	1004	0x17761453		
addi \$t1, \$t2, 0x01	1004		4	
sll \$t1,\$t2,2	1004		12	
clear \$t1	1004		0	
jal \$20	20	1004		
bge \$t1,\$t2,40	40			
ble \$t1,\$t2,40	1008			

see problem 3.10 pseudoinstructions.

slt \$at,\$t2,\$t1

beq \$at. \$zero. 40

Problem 3. (25%) Translate the following C code into MIPS. Please comment your code.
 Assume **x** is \$s0; **y** is \$s1; **p** is \$a0 and points to integers; **s** is \$a1 and points to unsigned char.
Points will be taken off for assembler syntax errors.

(a) **x = (x + (y - 3)) & 0xFF;**

```

addi      $t0, $s1, -3
add      $t0, $s0, $t1
andi      $s0, $t0, 0xFF

```

(b) **x = (y > 3)? x + 1 : x - 1;**

```

addi      $t0, $zero, 3
slt       $t1, $t0, $s1          # if ((3 < y )== 0) goto L1
beq      $t1, $zero, L1
addi      $s0, $s0, 1           # then
j        L2
L1: addi      $s0, $s0, -1      # else
L2:

```

(c) **for(x=1; x; x++) { y++; }**

```

addi      $s0, $zero, 1          # x = 1;
L2: beq      $s0, $zero, L1      # if (x == 0) goto L1
addi      $s1, $s1, 1           # y++ same as y=y+1
addi      $s0, $s0, 1           # x++
j        L2;
L1:

```

(d) **p[5] = *p + y + *s;** /*note: same as *(p+5) = *(p+0) + y + *(s + 0) */

```

lbu      $t0,0($a1)      # *s      (points to unsigned char)
add      $t1,$t0,$s1      # + y
lw       $t2,0($a0)      # *p      (points to int)
add      $t3, $t2, $t1    # *p + y + *s
addi     $t4,$zero,20    # 5 * sizeof(int) = 5*4 = 20
add      $t5,$t4,$a0      # address of p[5]
sw       $t3,0($t5)      # *(p+5) =

```

(e) **p[x+3] = s[x+2] | 0x01** /* note: same as *(p+x+3) = *(s+x+2) | 0x01 */

```

addi     $t0,$s0,2          # x + 2
add      $t1,$a1,$t0        # (s + x + 2)
lbu      $t2,0($t1)        # = *(s + x + 2)
ori      $t3,$t2,0x01
addi     $t4,$s0,3          # x+3
add      $t5,$t4,$t4        # 2*(x+2)
add      $t6,$t5,$t5        # 4*(x+2)
add      $t7,$a0,$t6        # p + x + 2
sw       $t3,0($t7)        # *(p+x+2) = $t3

```

alternate solution
addi \$t0,\$s0,2
add \$t1,\$a1,\$t0
lbu \$t2,0(\$t1)
ori \$t3,\$t2,0x01
addi \$t4,\$s0,3
add \$t5,\$t4,\$s0
add \$t6,\$t5,\$t5
add \$t7,\$a0,\$t6
sw \$t3,12(\$t7)

Problem 4. (25%) Translate the following code and add comments

```

void rgcd(int x, int y)  {
    register int w=0, z=1;

    if (y == 0) { return x; }
    z = x + y;
    w = rgcd(y, z);
    return w + z;
}

```

(a) Write the prolog

```

.text                      # read appendix A.6 and A
rgcd:
    add $sp,$sp -12          # allocate space
    sw  $ra,8($sp)
    sw  $s1,4($sp)
    addi $s1,$zero,1          # z = 1;
    sw  $s0,0($sp)
    add $s0,$zero,$zero        # w = 0

```

(b) Write the body

```

bne $a1,$zero,rgcd2      # if (y != 0) goto rgcd2
add $v0,$a0,$zero          # return x
j   rgcd_epilog

rgcd2:
    add $s1,$a0,$a1          # z=x + y
    add $a0,$a1,$zero          # $a0 = y
    add $a1,$s1,$zero          # $a1 = z
    jal rgcd
    add $s0,$v0,$zero          # w = return
    add $v0,$s0,$s1            # return = w+z

```

(c) Write the epilog

```

rgcd_epilog:
    lw   $s0,0($sp)           # not require to restore $a0, $a1
    lw   $s1,4($sp)           # restore caller's $s0
    lw   $ra,8($sp)            # restore caller's $s1
    addi $sp,$sp,12             #free space
    jr  $ra

```

Problem 5. (10%) Translate the following global variables

(a) char x[6] = “hello”;

```
.data
x:      .byte    'h','e','l','l','o',0
```

(b) struct treenode {
 struct treenode *left, *right;
 char symbol[4];
} root;

```
.data
root:
root_left:     .word    0
root_right:    .word    0
root_symbol:   .byte    0,0,0,0
```

Problem 6. (15%) Given the following instruction sequence in the table below.

Assume the (alu instructions are 6 clocks); (loads 8 clocks); (stores 7 clocks); (jumps 2 clocks); (branches 5 fall through/8 for branch);

- (a) Show the **best** case timing path through the code showing annotations and total.
- (b) Show the **worst** case timing path through the code showing annotations and total.
- (c) What values will make this code execute the worst case?

\$s0 = 0 and \$s1 = 3

instruction	best case	worst case
bne \$s0,\$zero,L2	8 branch	5 fall through 5 fall
addi \$t0,\$zero,3		6 alu 6 alu
bne \$s1,\$t0,L2		5 fall through 8 branch
addi \$s1,\$s1,1		6 alu
j L1		2 jump
L2: addi \$s1,\$zero,10	6 alu	6 alu
L1:		
Time	14 clocks	24 clocks 25 clocks ¹ (correct)

¹correct answer due to a quick thinking student