

EECS 281:
Name: _____

Test 3 (5 pages)
Email: _____

Due: Tuesday, November 2, 2004
Grade: _____ (100 points max)

1. (10 points) Please answer the following True or False in the context of Boolean Algebra:

$$\textcircled{T} \quad F \quad \bar{b} + a = \bar{a}\bar{b} + a \quad \bar{b} + a = \bar{b}(\bar{a} + \bar{a}) + \bar{a} = \bar{b}a + \bar{a}\bar{b} + a = a(\bar{b} + \bar{a}) + \bar{a}\bar{b} \\ = a + \bar{a}\bar{b} = \text{RHS}$$

$$\textcircled{T} \quad F \quad \sum_{ab}(1, 2) = \bar{a} \oplus \bar{b} \quad \bar{a} \oplus \bar{b} = \bar{a}\bar{b} + \bar{b}\bar{a} = \bar{a}\bar{b} + \bar{a}\bar{b} = \Sigma_{ab}(1, 2)$$

$$T \quad \textcircled{F} \quad \sum_{abc}(1, 8, 4, 2, 11, 13, 14, 7) = \prod_{abc}(3, 5, 6, 9, 10, 12, 0) \quad \text{RHS} \not\equiv \sum_{abcd} = \pi_{abcd}(3, 5, 6, 9, 10, 12, 0, 15)$$

$$\textcircled{T} \quad F \quad \prod_{abc}(0, 4, 7) = (\bar{a} + \bar{b} + \bar{c})(a + b + c)(\bar{a} + b + c)$$

$$T \quad \textcircled{F} \quad \bar{a}(b + c) = \bar{a} + \bar{a}b + \bar{a}c \quad LHS = \bar{a}b + \bar{a}c \quad \text{RHS} = \bar{a} + \bar{a}b + \bar{a}c \quad \text{cannot be simplified}$$

2. (10 points) Use Boolean Algebra to establish the identity. Show the Theorem numbers (i.e. T1-T13) for each step of your proof:

| Theorem | Expression |
|---------|---|
| | $a = (\bar{a} + \bar{c})(\bar{a} + c) + (a + b)(a + c)a$ |
| | $RHS = (\bar{a} + \bar{c})(\bar{a} + c) + (a + b)(a + c)a$ |
| | $= (\bar{a} + \bar{c}) + (\bar{a} + c) + (a + b)(a + c)a$ |
| | $= (\bar{a} \cdot \bar{c}) + (\bar{a} \cdot c) + (a + b)(a + c)a$ |
| | $= a \cdot c + a \cdot \bar{c} + (a + b)(a + c)a$ |
| | $= acc + \bar{c} + a(a + b)(a + c)$ |
| | $= a + a(a + b)(a + c)$ |
| | $= a[1 + (a + b)(a + c)]$ |
| | $= a[1]$ |
| | $= a = LHS$ |
| | |
| | |
| | |
| | |
| | |

$$\begin{aligned}
 3a) \quad \bar{a}\bar{c}\bar{d} &= a\bar{c}\bar{d}(b+\bar{b}) = ab\bar{c}\bar{d} + a\bar{b}\bar{c}\bar{d} \\
 \bar{c}\oplus d &= \bar{c}\bar{d} + cd = \bar{c}\bar{d}(b+\bar{b}) + cd(b+\bar{b}) \\
 &= b\bar{c}\bar{d} + \bar{b}\bar{c}\bar{d} + bcd + \bar{b}cd \\
 &= b\bar{c}\bar{d}(a+\bar{a}) + \bar{b}\bar{c}\bar{d}(a+\bar{a}) + bcd(a+\bar{a}) \\
 &\quad + \bar{b}cd(a+\bar{a}) \\
 &= ab\bar{c}\bar{d} + \bar{a}b\bar{c}\bar{d} + a\bar{b}\bar{c}\bar{d} + \bar{a}\bar{b}\bar{c}\bar{d} + abcd \\
 &\quad + \bar{a}bcd + abcd + \bar{a}\bar{b}cd
 \end{aligned}$$

$$\begin{aligned}
 \bar{b} &= \bar{b}(a+\bar{a}) \\
 &= a\bar{b} + \bar{a}\bar{b} \\
 &= \bar{a}\bar{b}(c+c) + \bar{a}\bar{b}(cc+\bar{c}) \\
 &= \bar{a}\bar{b}c + \bar{a}\bar{b}\bar{c} + \bar{a}\bar{b}c + \bar{a}\bar{b}\bar{c} \\
 &= \bar{a}\bar{b}c[d+\bar{d}] + \bar{a}\bar{b}\bar{c}[d+\bar{d}] + \bar{a}\bar{b}c[d+\bar{d}] \\
 &\quad + \bar{a}\bar{b}\bar{c}[d+\bar{d}] \\
 &= a\bar{b}cd + a\bar{b}c\bar{d} + \bar{a}\bar{b}\bar{c}d + a\bar{b}\bar{c}\bar{d} + \bar{a}\bar{b}cd + \bar{a}\bar{b}\bar{c}\bar{d} \\
 &\quad + \bar{a}\bar{b}\bar{c}d + \bar{a}\bar{b}\bar{c}\bar{d}
 \end{aligned}$$

$$f(a,b,c,d) \in \{0, 1, 2, 3, 4, 7, 8, 9, 10, 11, 12, 15\}$$

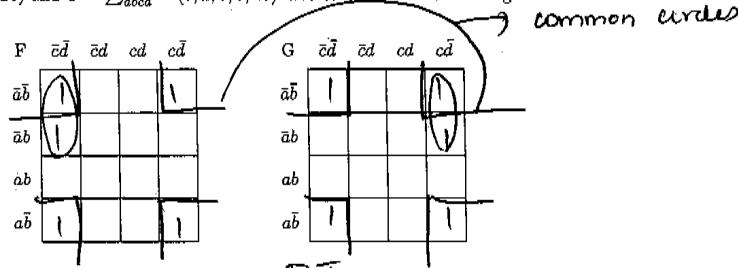
3a. (15 points) Show the optimal minimal circling in the k-map in minterm function $f(a, b, c, d) = a\bar{c}\bar{d} + (\bar{c} \oplus d) + \bar{b}$:

| | $\bar{c}\bar{d}$ | $\bar{c}d$ | cd | $c\bar{d}$ |
|------------------|------------------|------------|------|------------|
| $\bar{a}\bar{b}$ | 1 | 1 | 1 | 1 |
| $\bar{a}b$ | 1 | | 1 | |
| $a\bar{b}$ | 1 | | 1 | |
| ab | 1 | 1 | 1 | 1 |

3b. Give the $\sum_{abcd} = \underline{\underline{0, 1, 2, 3, 4, 7, 8, 9, 10, 11, 12, 15}}$

3c. Give MSOP = $\underline{\underline{b + \bar{c}\bar{d} + cd = \bar{b} + c\bar{c} \oplus d}}$

4a. (20 points) Show the optimal multi-output minimal circling the terms and in the k-map in minterm function $F = \sum_{abcd} = (0, 2, 4, 8, 10)$ and $G = \sum_{abcd} = (0, 2, 6, 8, 10)$. Indicate which circle belongs to what function.

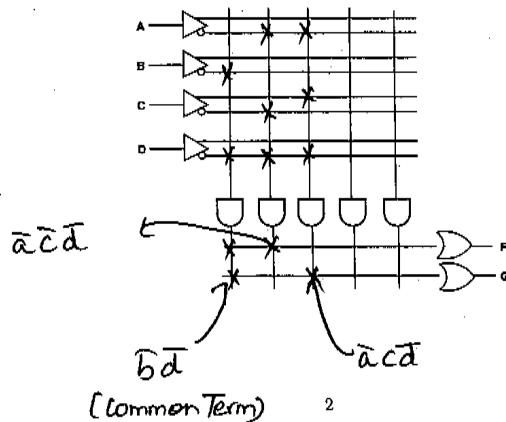


4b. Give the common term of multi-output MSOP = $\underline{\underline{b\bar{d}}}$

4c. Give the multi-output MSOP of $F = \underline{\underline{b\bar{d} + \bar{a}\bar{c}\bar{d}}}$

4d. Give the multi-output MSOP of $G = \underline{\underline{b\bar{d} + \bar{a}c\bar{d}}}$

4e. Fill in the PLA



$$\Sigma(a, b, c, d) \ (0, 15, 15, 10, 12, 8)$$

| | | | | | | |
|---|--|---|--|---|---|---|
| $\begin{array}{r} 60 \\ 61 \\ 62 \\ 63 \\ 64 \end{array}$ | $\begin{array}{r} \sqrt{0} \\ \sqrt{2} \\ \sqrt{8} \\ 5 \\ 15 \end{array}$ | $\begin{array}{r} 0000 \\ 0010 \\ 1000 \\ 0101 \\ 1010 \\ 1111 \end{array}$ | $\begin{array}{r} (0, 12) \\ (0, 8) \\ (2, 10) \\ (8, 10) \end{array}$ | $\begin{array}{r} 00-0 \\ -000 \\ -010 \\ 10-0 \end{array}$ | $\begin{array}{r} (0, 12, 8, 10) \\ (0, 8, 2, 10) \\ (0, 2, 5, 10) \end{array}$ | $\begin{array}{r} -0-0 \\ -0-0 \\ -0-0 \end{array}$ |
|---|--|---|--|---|---|---|

$f = \bar{b}\bar{d} + abcd + \bar{a}b\bar{c}d$

5a. (20 points) Do the Quine-McCluskey Algorithm of $\sum_{a,b,c,d}(0, 5, 15, 10, 2, 8)$.

| Group | Minterms | 0-cubes | Minterms | 1-cubes | Minterms | 2-cubes |
|-------|----------|--------------|----------------|--------------|----------------------|--------------|
| G_0 | 10 | 0000 | 012 18 | 00-0 -000 | 0128, 10 1218, 10 | -0-0 -0-0 |
| G_1 | 12 18 | 0010 1000 | 2, 10 8, 10 | -010 10-0 | | |
| G_2 | 5 | 0101 | | | | |
| G_3 | 10 | 1010 | | | | |
| G_4 | 15 | | | | | |

5b. Fill in the covering table

| EPI? | Needed? | PI-cubes | 0 | 5 | 15 | 10 | 2 | 8 |
|------|---------|----------|---|---|----|----|---|---|
| ✓ | Yes | 0101 | ✓ | | | | | |
| ✓ | Yes | 1111 | | | ✓ | | | |
| ✓ | Yes | -0-0 | ✓ | | ✓ | ✓ | ✓ | |
| | | | | | | | | |
| | | Covered? | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ |

5c. Give the MSOP = $b\bar{d} + abc\bar{d} + \bar{a}b\bar{c}\bar{d}$

5d. Show the optimal k-map:

| | $\bar{c}\bar{d}$ | $\bar{c}d$ | cd | $c\bar{d}$ |
|------------------|------------------|------------|------|------------|
| $\bar{a}\bar{b}$ | 1 | | 1 | |
| $\bar{a}b$ | | 1 | | |
| $a\bar{b}$ | | | 1 | |
| ab | 1 | | | 1 |

5e. Give the MSOP of the k-map: $\bar{b}\bar{d} + \bar{a}b\bar{c}\bar{d} + abcd$

$$a^n b^m c$$
$$=((a \oplus b) \oplus c)$$

$$= (\bar{a}\bar{b} + \bar{a}b) \oplus c$$
$$= \overline{(\bar{a}\bar{b} + \bar{a}b)} c + \bar{c} (a\bar{b} + \bar{a}b)$$
$$= [(\bar{a} + \bar{b}) \cdot (\bar{a} + b)] c + a\bar{b}\bar{c} + \bar{a}b\bar{c}$$
$$= [(\bar{a} + b) c \cdot (a + \bar{b}) c] + [a\bar{b}\bar{c} + \bar{a}b\bar{c}]$$
$$= [(\bar{a}c + bc) \cdot (ac + \bar{b}c)] + [a\bar{b}\bar{c} + \bar{a}b\bar{c}]$$
$$= abc + \bar{a}\bar{b}c + a\bar{b}\bar{c} + \bar{a}b\bar{c}$$

6a. (10 points) Given $\sum_{a,b,c,d}(0, 5, 15, 10, 2, 8)$ and the don't cares (7, 13), show the optimal k-map:

| | $\bar{c}\bar{d}$ | $\bar{c}d$ | cd | $c\bar{d}$ |
|------------------|------------------|------------|------|------------|
| $\bar{a}\bar{b}$ | 1 | | | |
| $\bar{a}b$ | | 1 | d | |
| $a\bar{b}$ | | a | 1 | |
| ab | 1 | | | |

6b. Give the MSOP of the k-map: $\bar{b}\bar{d} + b\bar{d} = (\bar{b} \oplus d)$

7. (15 points) A programmer has written the following C code fragment:

```
f=1;
if (~a) {
    if (~(b ^ c)) { f=0; }
}
else if (b ^ c) { f=0; }
```

7a. Give the truth table for the variable f (assume that a, b, c are boolean values only):

| a | b | c | f | $\Sigma_{abc} f = (1, 2, 4, 7)$ |
|---|---|---|---|---------------------------------|
| 0 | 0 | 0 | 1 | |
| 0 | 0 | 1 | | |
| 0 | 1 | 0 | | |
| 0 | 1 | 1 | 0 | |
| 1 | 0 | 0 | | |
| 1 | 0 | 1 | 0 | |
| 1 | 1 | 0 | 0 | |
| 1 | 1 | 1 | 1 | |

7b. Give the optimal k-map of 7a.

| | $\bar{b}\bar{c}$ | $\bar{b}c$ | bc | $b\bar{c}$ |
|-----------|------------------|------------|------|------------|
| \bar{a} | | 1 | | 1 |
| a | 1 | | 1 | |

7c. Give the MSOP of the k-map: $\bar{a}\bar{b}\bar{c} + \bar{a}\bar{b}c + a\bar{b}\bar{c} + abc$ [cannot be further minimized from]

7d. Re-write as optimal C code:
 $f = ((\bar{a} \& \bar{b} \& \bar{c}) | (\bar{a} \& b \& \bar{c}) | (a \& \bar{b} \& \bar{c}) | (a \& b \& c))$ [K-Map]

or the best solution would be $f = (a \wedge b \wedge c)$

7e. (5 points extra credit) Re-write as optimal C code using minimal parenthesis and as many exclusive-or's as possible:

$$f = a \wedge b \wedge c$$

x1. (10 points extra credit) Using C++ data types for a machine that uses a char of 5-bits, convert the following into one's complement big-endian binary and if not, then show why not?: where signed char s, a=-1, b=6; For addition and indicate if end-around-carry, overflow and/or carry has occurred. Show work.

| | |
|---------------------------|--|
| Give unsigned char range: | $0 \text{ to } 2^5 \Rightarrow 0 \text{ to } 32$ |
| Give signed char range: | $-(2^{n-1}-1) \text{ to } (2^{n-1}-1) = -15 \text{ to } +15$ |
| unsigned char x = 30; | 11110 |
| signed char x = -1; | -1 (\Rightarrow) 11110 |
| s = (~a)+1; | $\Rightarrow 00010$ |
| s = ~a; | $\Rightarrow 00001$ |
| s = -b; | $-b \Rightarrow 11001$ |
| s = a & b; | $\Rightarrow 00110$ |
| s = a + b; | $\Rightarrow 00101 \Rightarrow +5$ |
| s = a - b; | $-7 \Rightarrow -1-6 = 11000$ |

| Theorem | Relationship | Dual | XOR | Property |
|---------|---|------------------------------------|---|-----------------------------|
| T1 | $a1 = a$ | $a+0 = a$ | $a \oplus 0 = a$ | Identity |
| T2 | $a0 = 0$ | $a+1 = 1$ | $a \oplus 1 = \bar{a}$ | Domination |
| T3 | $aa = a$ | $a+a = a$ | $a \oplus a = 0$ $a \oplus a \oplus a = a$ | Idempotency |
| T4 | $\bar{\bar{a}} = a$ | | | Involution |
| T5 | $a\bar{a} = 0$ | $a+\bar{a} = 1$ | $a \oplus \bar{a} = 1$ | Complement |
| T6 | $ab = ba$ | $a+b = b+a$ | $a \oplus b = b \oplus a$ | Commutative |
| T7 | $(ab)c = a(bc)$ | $(a+b)+c = a+(b+c)$ | $(a \oplus b) \oplus c = a \oplus (b \oplus c)$ | Associative |
| T8 | $(a+b)(a+c) = a+bc$ | $a(b+c) = ab+ac$ | $a(b \oplus c) = ab \oplus ac$ | Distributive |
| T9 | $a(a+b) = a$ | $a+ab = a$ | $a \oplus ab = \bar{a}$ | Absorption Covering |
| T10 | $(a+b)(a+\bar{b}) = a$ | $ab+a\bar{b} = a$ | $ab \oplus a\bar{b} = a$ | Combining |
| T11 | $(a+b)(\bar{a}+c)(b+c) = (a+b)(\bar{a}+c)$ | $ab+\bar{a}c+bc = ab+\bar{a}c$ | | Consensus Proof by k-map |
| T12 | $a+a+\dots+a = a$ | $aa\dots a = a$ | $a \oplus a \oplus \dots \oplus a_{odd} = a$ $a \oplus a \oplus \dots \oplus a_{even} = 0$ | Generalized Idempotency |
| T13 | $\overline{a+b} = \bar{a}\bar{b}$ | $\bar{a}\bar{b} = \bar{a}+\bar{b}$ | $\bar{a}\bar{b} = \bar{a} \oplus \bar{b} \oplus \bar{a}\bar{b}$ | DeMorgan |
| XOR | $ab = a \oplus \bar{b} \oplus \bar{a}\bar{b}$ | $a+b = a \oplus b \oplus ab$ | $a \oplus b = \bar{a} \oplus \bar{b} = a\bar{b} + \bar{a}b$ | Definition |